

# Automatic selection of contiguous forest reserves

SSAFR

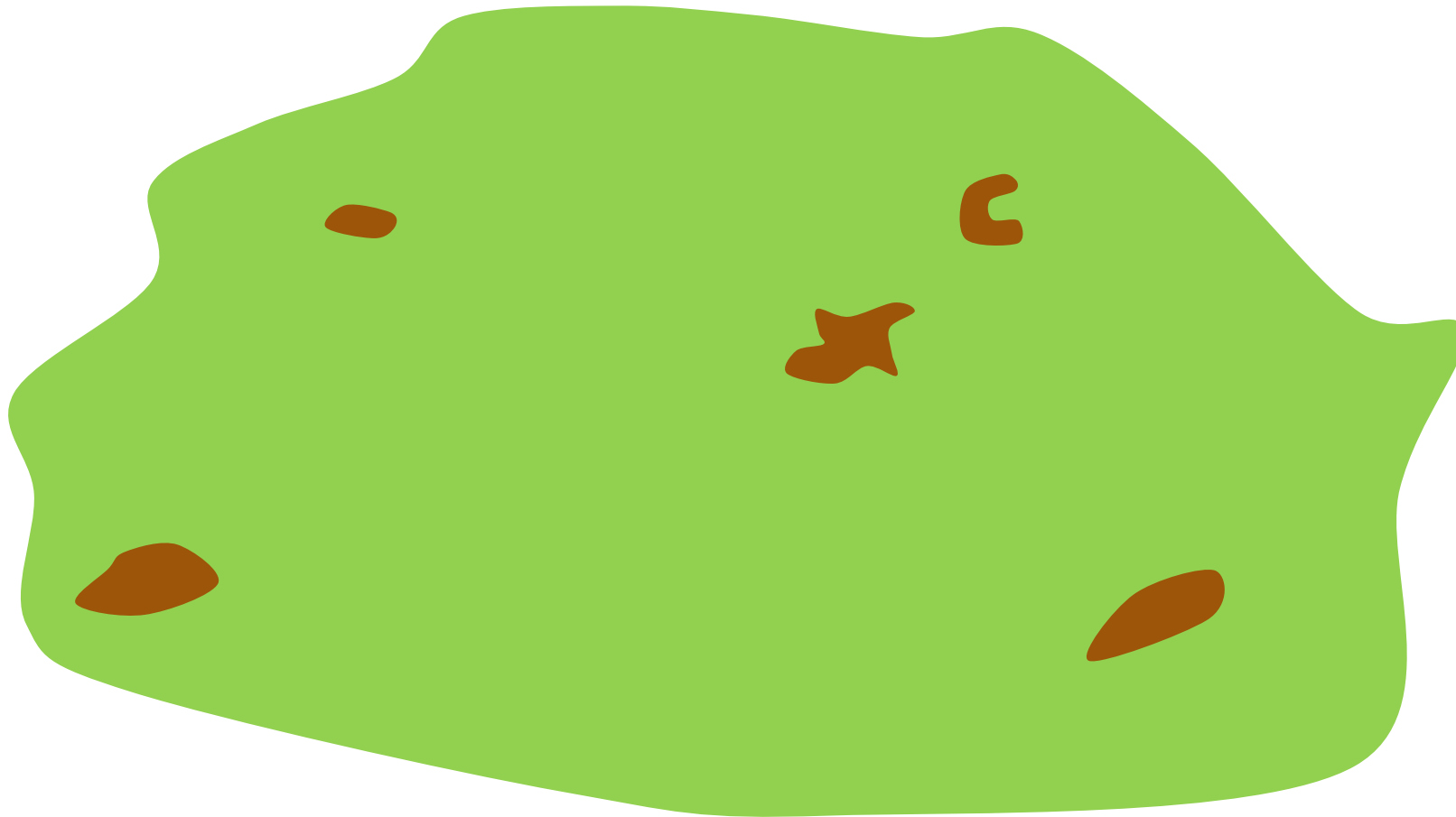
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# Problem – selection of reserves



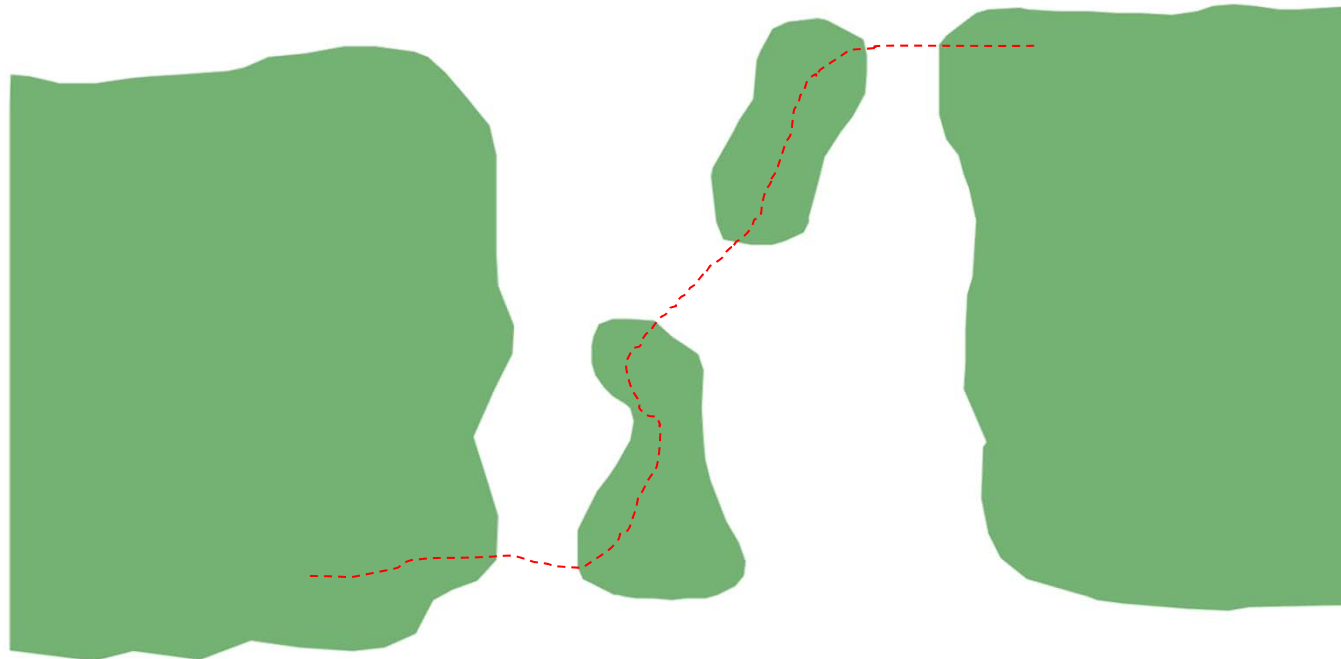
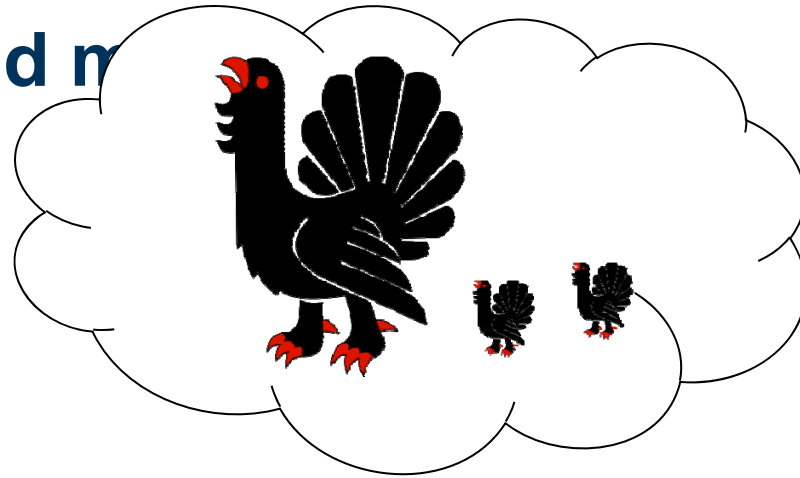
# Problem – real world situation

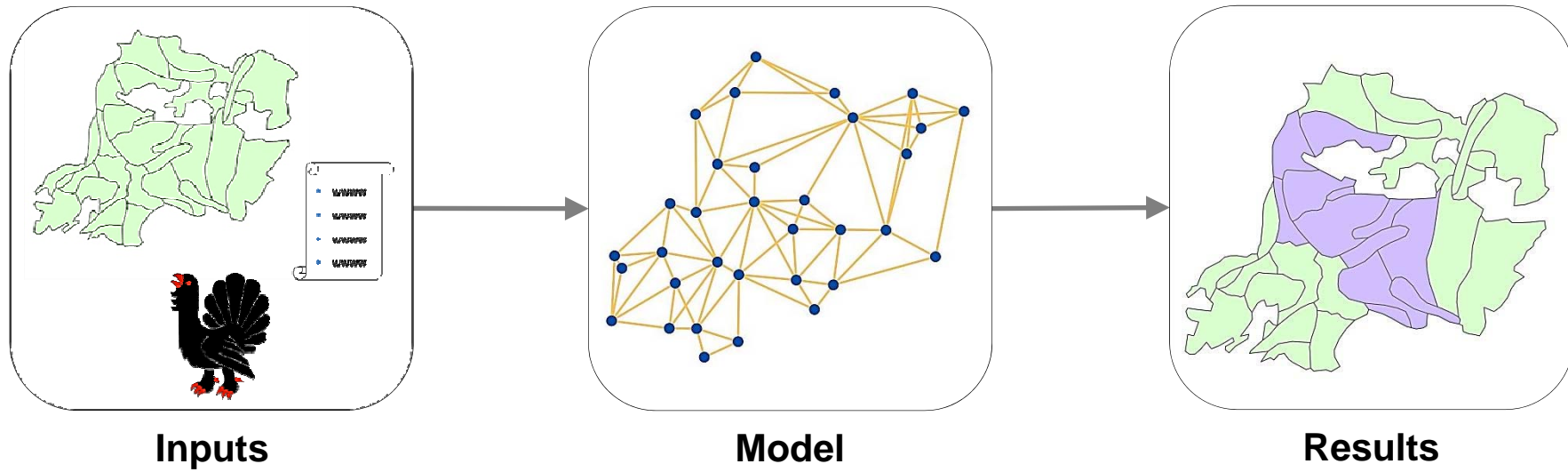


## Aim of the thesis

Develop a general model for an automatic selection of contiguous forest reserves considering ecological and economical aspects

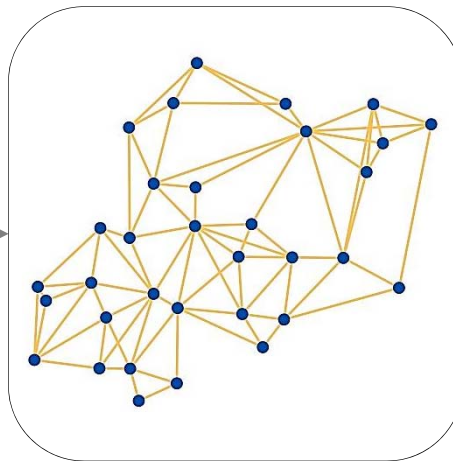
# Continuity and m



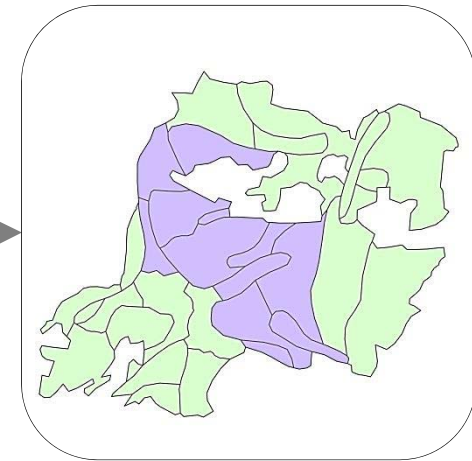




**Inputs**



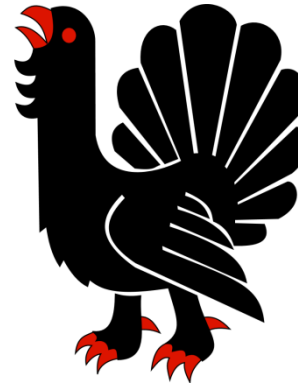
**Model**



**Results**

# Inputs

- **Conservation aim**
- Decision units
- Constraints
- Assessment criteria incl. weighting





# Input

- Conservation aim
- **Decision units**
- Constraints
- Assessment criteria  
incl. weighting



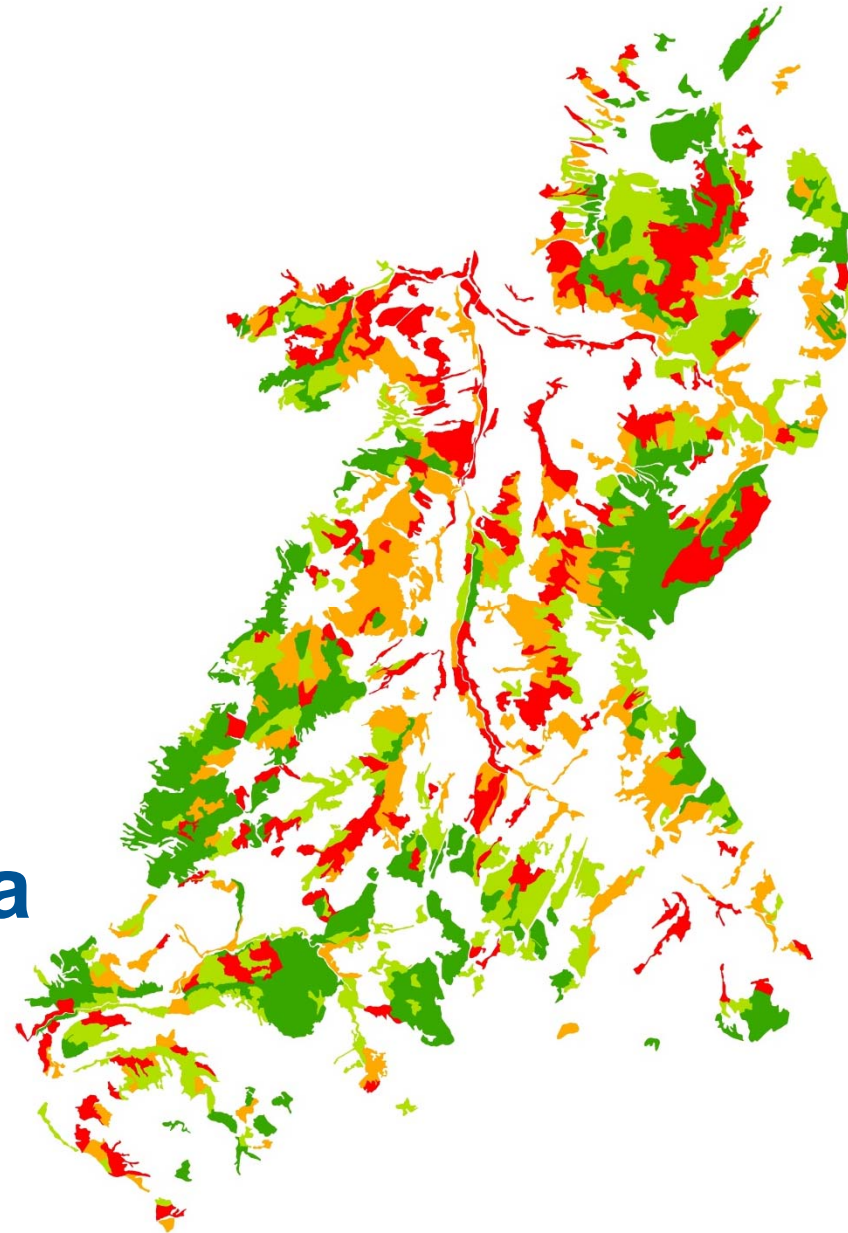
# Input

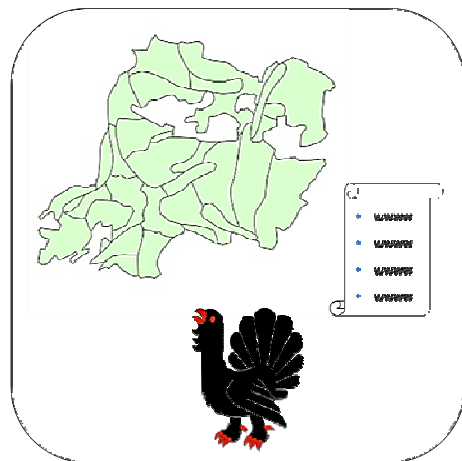
- Conservation aim
- Decision units
- **Constraints**
- Assessment criteria incl. weighting



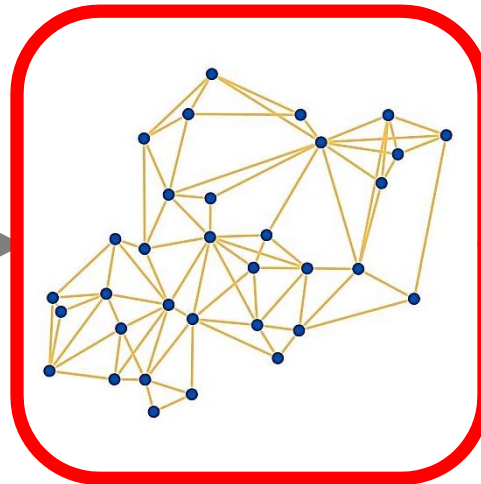
## Input

- Conservation aim
- Decision units
- Constraints
- **Assessment criteria  
incl. weighting**

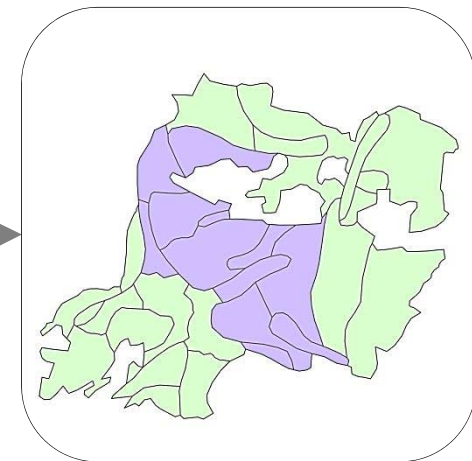




**Inputs**



**Model**



**Results**

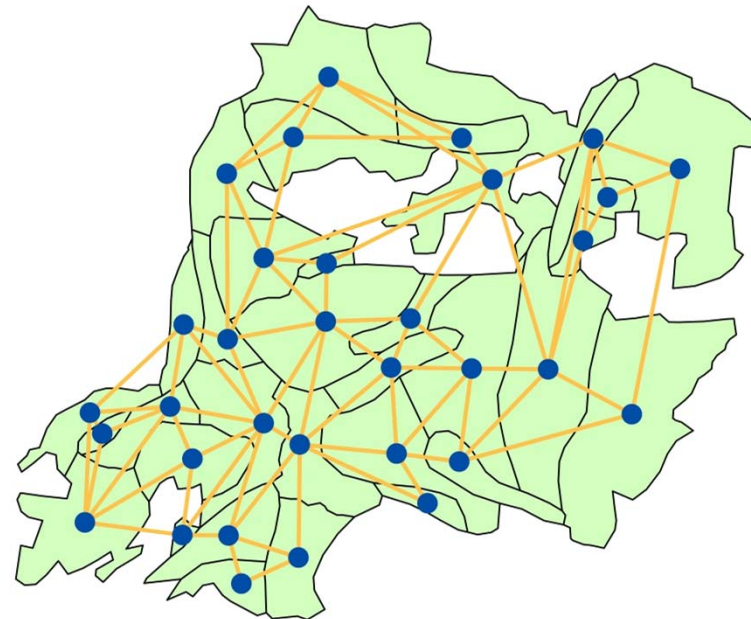
# Optimization

## Integer Linear Programming (ILP)

**Objective:** Maximize the suitability of the selected patches

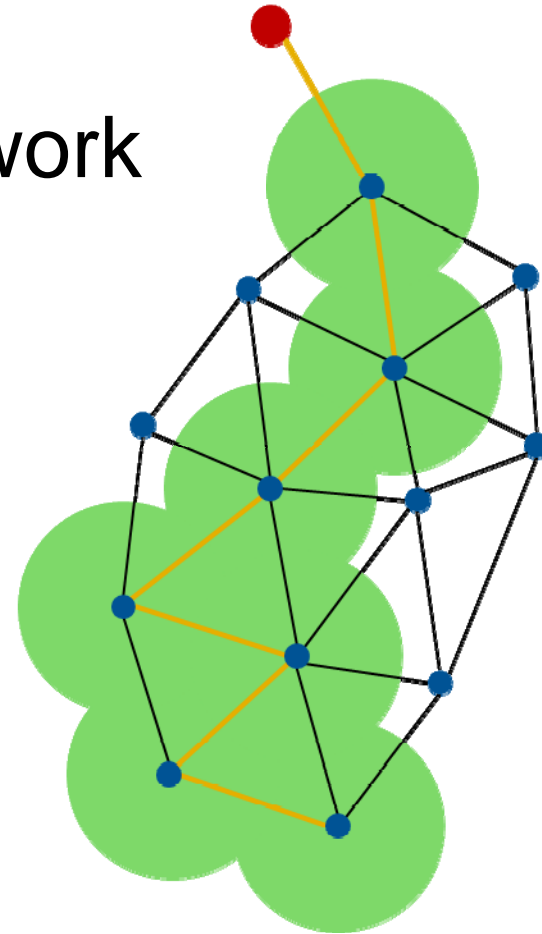
### Constraints:

- Maximum area
- Number of regions
- Contiguity



# ILP

- Forest represented as a network
- Add exit node
- Search a path from start node to exit node



# Model parameters

letter	description
$x_i$	1 if patch $i$ is selected, 0 otherwise
$y_{i0}$	1 if edge between patch $i$ and exit node is selected, 0 otherwise
$y_{ij}$	1 if edge between patch $i$ and $j$ is selected, 0 otherwise
$y_{ji}$	1 if edge between patch $j$ and $i$ is selected, 0 otherwise
$h_i$	Suitability of patch $i$
$B$	Maximum area
$c_i$	Area of patch $i$
$R$	Maximum quantity of regions
$I$	set of patches $i$
$N$	Number of patches
$N_i$	Set of adjacent patches $j$ of patch $i$
$G$	Set of patches $i$ of the cycle (=potential start node)

# Analytical Model (ILP)

**Objective Function**

$$\text{Max} \sum_{i=1}^N \sum_{j \in N_i} y_{ji} * h_i$$

*Maximize suitability*

**s.t.**

$$\sum_{i=1}^N x_i * c_i \leq B$$

*Maximum area*

$$\sum_{i=1}^N y_{i0} \leq R$$

*Number of regions*

$$\sum_{j \in N_i} y_{ij} \geq x_i,$$

for all  $i \in I$

$$y_{ji} \leq \sum_{j \in N_i} y_{ij}$$

for all  $i \in I$

$$y_{ij} \leq x_i,$$

for all  $i \in I$

$$y_{ij} + y_{ji} \leq 1$$

$$x_i \in \{0,1\}, y_{ij} \in \{0,1\}$$

*Binarity*

*Continuous  
flow*

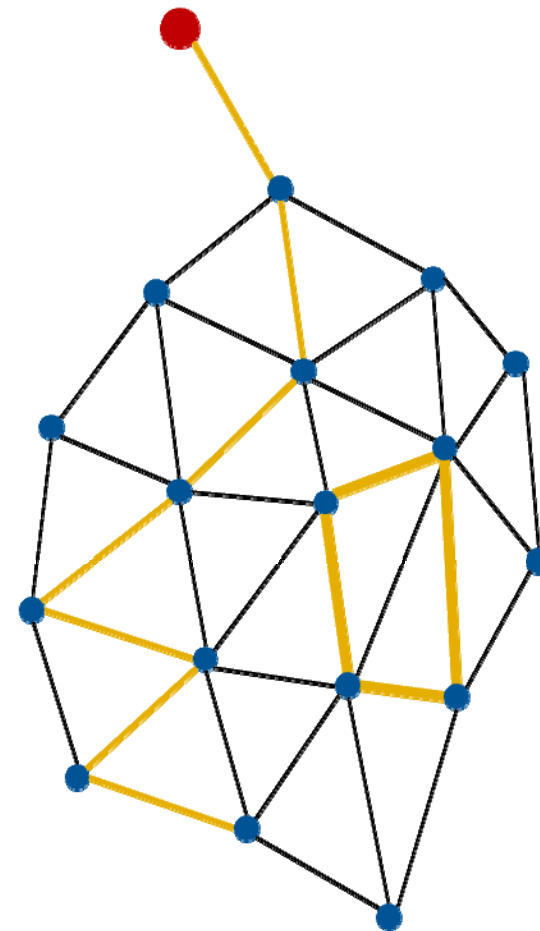


# Cycle prohibiting constraints

Solutions with several cycles

**Introduce new constraints  
which prohibit cycles**

(from Church and Cova, 2000)



# Cycle prohibiting constraints

**u** is the equivalent of **y** in the cycle

**w<sub>i</sub>**: 1 if node *i* is the start node of the cycle, 0 otherwise

**e**: 1 if a cycle exists, 0 otherwise

$$x_i \leq e,$$

for all  $i \in G$

$$e \leq \sum_{i \in G} w_i$$

$$w_i + \sum_{j \in N_i} u_{ji} = \sum_{j \in N_i} u_{ij},$$

for all  $i \in G$

$$\sum_{j \in N_i} u_{ji} = \sum_{j \in N_i} u_{ij},$$

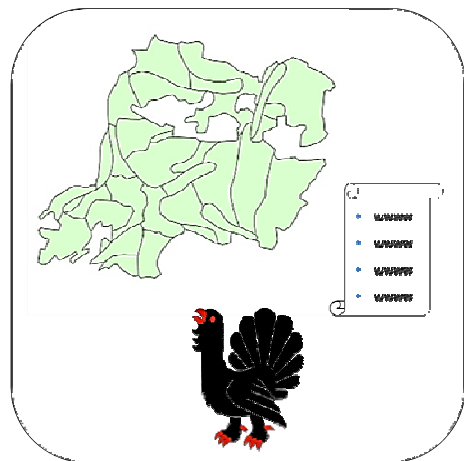
for all  $i \notin G$

$$e = \sum_{i=1}^N u_{i0}$$

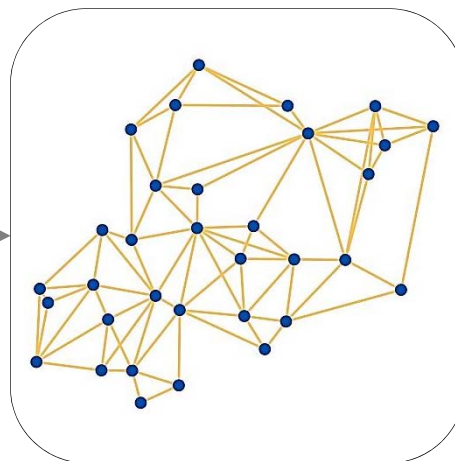
$$u_{ij} \leq y_{ij},$$

for all  $i \in I$

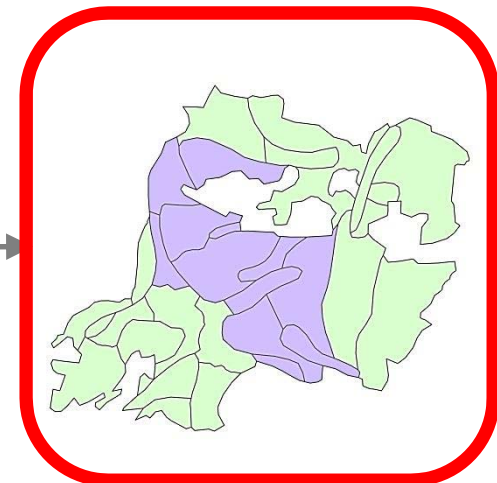
$$e \in \{0,1\}, \quad u_{ij} \in \{0,1\}, \quad w_i \in \{0,1\}$$



**Inputs**

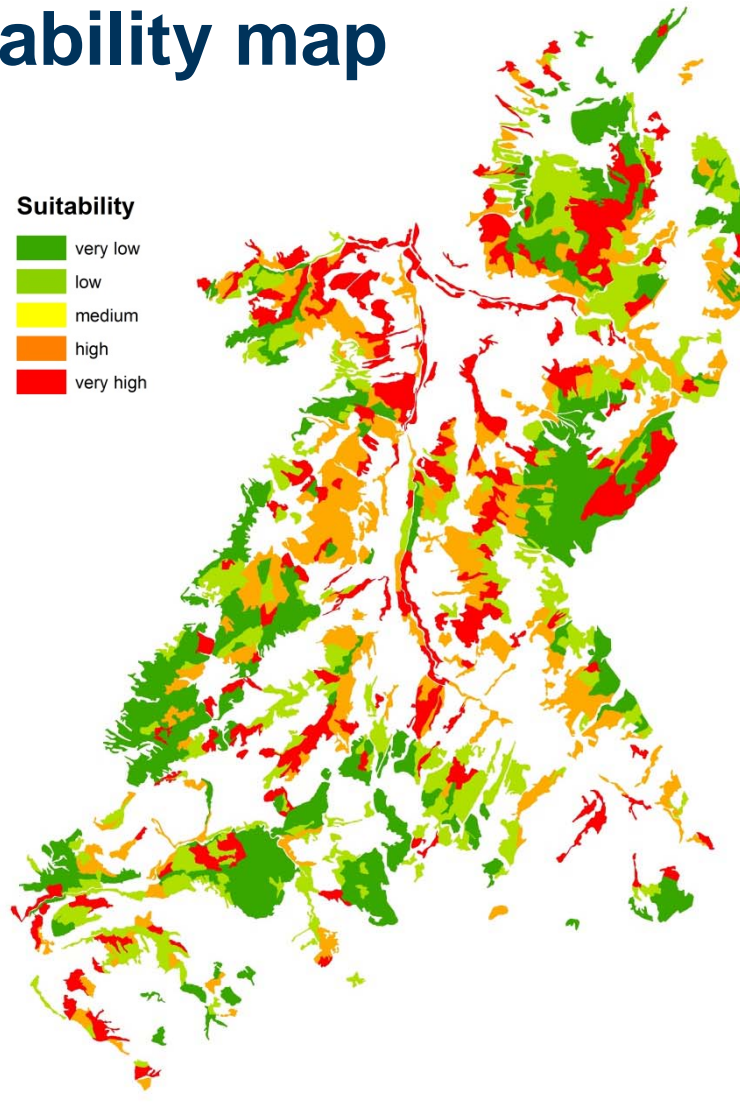


**Model**

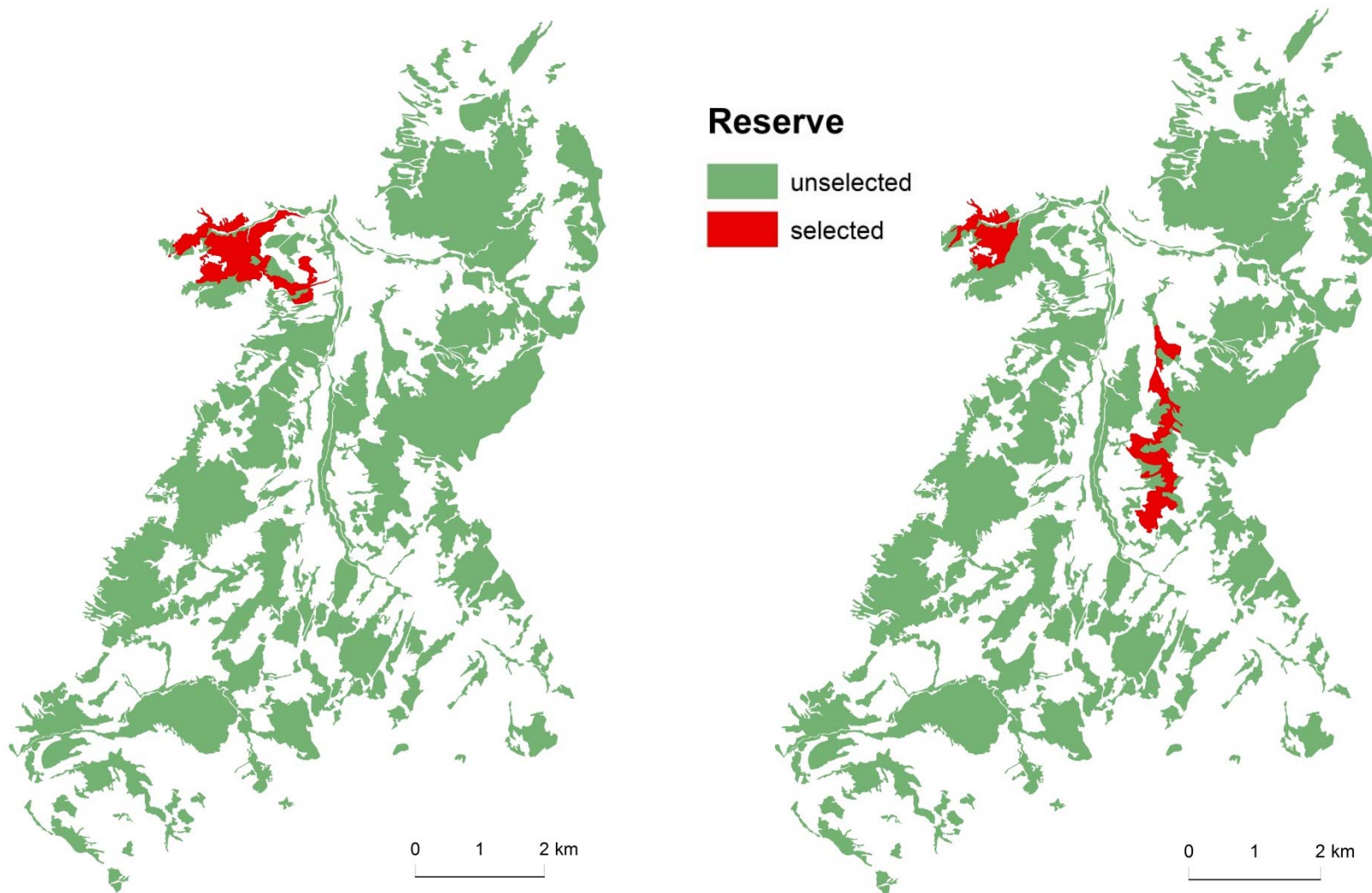


**Results**

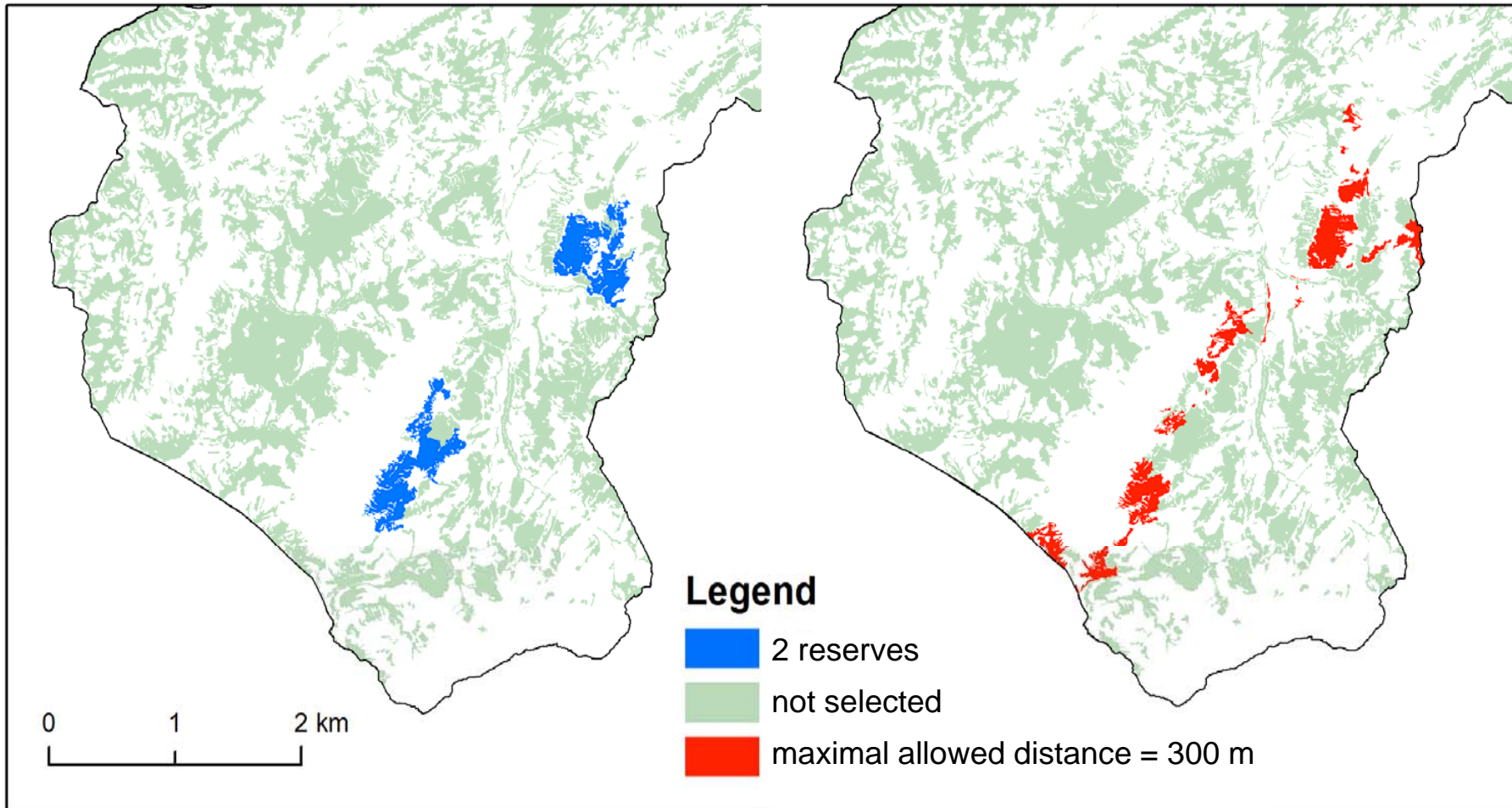
# Results – Suitability map



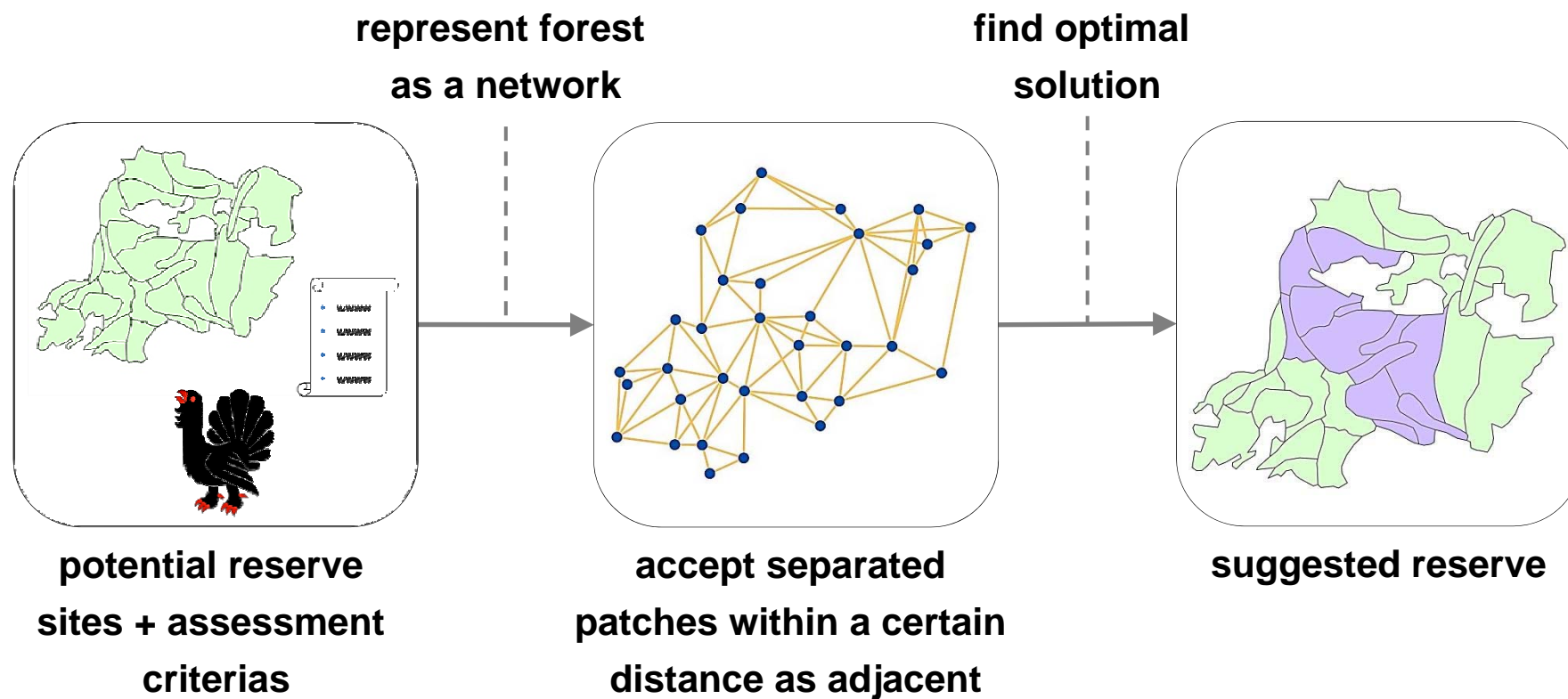
# Results – 1 region vs. 2 regions



## Results – stepping stones



# Summary



# Outlook

## Improvement

- Include barriers
- Other patch shapes
- Better geodata
- More computing power

## Need for research

- Better knowledge about behaviour and needs of species
- Faster mathematical models to ensure contiguity



A photograph of a vibrant green meadow filled with small white flowers in the foreground. The meadow leads to a dense forest of tall evergreen trees in the background. The sky is overcast and grey.

**Thanks a lot for your attention!**

# ILP formula part 1

Maximize the suitability

$$\text{Max} \sum_{i=1}^N \sum_{j \in N_i} y_{ji} * h_i$$

Budget constraint or maximum selected area must be adhered to:

$$\sum_{i=1}^N x_i * c_i \leq B$$

If patch  $i$  is selected, at least one link must lead away from it:

$$\sum_{j \in N_i} y_{ij} \geq x_i, \quad \text{for all } i \in I$$

If a link leads to patch  $i$ , at least one link must lead away from it:

$$y_{ji} \leq \sum_{j \in N_i} y_{ij} \quad \text{for all } i \in I$$

## ILP formula part 2

The number of links to the exit node, has to be less or equal than the allowed number of regions:

$$\sum_{i=1}^N y_{i0} \leq R$$

If a link leads to patch  $i$ , it must be selected:

$$y_{ij} \leq x_i, \quad \text{for all } i \in I$$

One link can only lead in one direction:

$$y_{ij} + y_{ji} \leq 1$$

The decision variables must be 0 or 1:

$$x_i \in \{0,1\}, y_{ij} \in \{0,1\}$$

# Cycle prohibiting constraints part 1

If there is a cycle, the contiguity constraints must be activated:

$$x_i \leq e, \quad \text{for all } i \in G$$

If the contiguity constraints are active, the flow must start from a node in the cycle:

$$e \leq \sum_{i \in G} w_i$$

Flow balance in the potential start nodes (Input = Output):

$$w_i + \sum_{j \in N_i} u_{ji} = \sum_{j \in N_i} u_{ij}, \quad \text{for all } i \in G$$

Flow balance in the nodes (except for start nodes):

$$\sum_{j \in N_i} u_{ji} = \sum_{j \in N_i} u_{ij}, \quad \text{for all } i \notin G$$

## Cycle prohibiting constraints part 2

If the contiguity network is to be activated, one link to the exit node must exist:

$$e = \sum_{i=1}^N u_{i0}$$

If the contiguity network is to be activated, a connection to the corresponding problem must exist:

$$u_{ij} \leq y_{ij}, \quad \text{for all } i \in I$$

The decision variables must be 0 or 1:

$$e \in \{0,1\}, \quad u_{ij} \in \{0,1\}, \quad w_i \in \{0,1\}$$