## Testing the efficiency of adjacency algorithms for solving unit restricted models

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## INTRODUCTION

Adjacency constraints – important part of spatial harvest scheduling models since GIS tools were developed



legal constraints, environmental constraints etc.

# Early solvers implementation limited number of variables and constraints

this can lead to a loss of efficiency in solving some problems (as discussed by Torres-Rojo and Brodie 1990).

the reduction of constraints less relevant (Crowe et al. 2003)

There are many possibilities how to add adjacency constraints into the model



Very simple to create



Pairwise constraints



High number of constraints



Constraints based on analytical algorithms (Yoshimoto, Brodie 1994)



Low number of constraints Mathematical modification is needed

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#### The effect of neighbourhood type



#### The goal of the work



- To compare time needed for creating and solving models with adjacency constraints:
  - based on 3 different analytical algorithms + pairwise constraints
  - for 3 different types of neighbourhood



Is it necessary to consume computer memory and time by using the analytical algorithms when the sophisticated softwares and powerful computers are available?



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#### **MATERIAL AND METHODS**



738 randomly generated spatial structures (1-738 polygons)



**2-3 optimization** were made for each spatial structure –2,000 optimization in total



Totally **12** variants = 4 types of adjacency constraints × 3 types of neighbourhood





Gurobi optimization software was used for analysis



Personal computer Intel® Core<sup>™</sup> i7-2600 CPU @ 3.40 GHz was used



The models were written in Java programming language



**Mathematical model** 

$$\max Z = \sum_{i=1}^{n} \sum_{p=1}^{5} v_{ip} \cdot x_{ip}$$
$$\sum_{p=1}^{5} x_{ip} \le 1 \quad \forall i = 1, 2, ..., n$$

Adjacency constraints



	Evaluated time				
	Type of adjacency constraints	Algorithm processing time	Constraints creating time	Solving time	
	Triangular adjacency matrix (TAM)				
$\begin{vmatrix} \mathbf{M} \cdot x \leq \mathbf{A} \cdot 1 \\ \mathbf{M} = \mathbf{A} + \mathbf{B} \end{vmatrix}$	Row adjacency matrix (RAM)				
	Triangular row adjacency matrix (TRAM)				
$x_i + x_j \le 1$	Pairwise constraints				

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### RESULTS

Noighboughood	Neighbours statistics				
Neignbournood	Mean	Standard deviation	Max / Min		
Neumann	3.9	2.9	28 / 0		
Moore	4.2	3.1	28 / 0		
Buffer	7.0	5.3	44 / 0		

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	Number of resulted constraints				
	ТАМ	RAM	TRAM	Pairwise constraints	
Neumann	0 - 522	0 – 423	0 - 423	0 – 1,472	
Moore	0 - 530	0-441	0 - 441	0 – 1,628	
Buffer	0 - 620	0 - 527	0 - 527	0 – 2,778	

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#### **PROCESSING TIME FOR DIFFERENT ALGORITHMS – Buffer**

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#### CONSTRAINTS CREATING TIME FOR DIFFERENT ALGORTIHMS – Buffer neighbourhoods



## **MODEL SOLVING TIME OF DIFFERENT ALGORITHMS – Buffer**

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#### **DISCUSSION AND CONCLUSION**

Time needed for creating constraints in the model and solving is strongly related to the number of neighbours



Solving time depends not only on the number of polygons but also on the spatial structure



Analytical algorithms reduce the total number of constraints, however they can lead to the loss of eficiency of branch-and-cut algorithm for solving the model

# Thank you for your attention!

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